# Efficient 2D-Navigation for Mobile Service Robots

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## Contribution

We propose a set of algorithms for efficiently performing Simultaneous Localization and Mapping (SLAM) as well as path planning and motion control. These allow a mobile service robot to efficiently navigate and accomplish complex tasks in dynamic indoor environments.

## SLAM

SLAM is addressed in terms of *incremental range image registration*. The first laser scan  $D_0$  makes up an initial point map  $M_0$ . Points from sub-

## Sensor: 2D Laser Scanner



## **Typical Applications**

"Move to the refridgerator!"



### Platform



sequent scans  $D_i$ , i > 0 are *matched* against the map  $M_{i-1}$ . The resulting transformation  $\mathbf{T}_i$  is used to correct the robot's pose and the position of the points, yielding the transformed point set  $\check{D}_i$ .  $\check{D}_i$  is then added to  $M_{i-1}$  forming  $M_i$  to account for new information in  $D_i$ . After registering N range scans  $D_{i,...,N}$ , the point map  $M_N$  contains all points measured in the robot's workspace:

> $M_N = \bigcup \{ \mathbf{\check{d}}_{i,j} \mid \mathbf{\check{d}}_{i,j} \in \check{D}_i \}.$ i = [0, N]

The key idea of *sparse point maps* is to avoid duplicate storage of points. A point  $\mathbf{\check{d}}_{i,j} \in \check{D}_i$  is not added to  $M_{i-1}$ , if the point-to-point distance to its closest point  $\mathbf{m}_{i-1,k} \in M_{i-1}$  is smaller than a minimum allowable distance  $\epsilon_{\rm D}$ :

$$M_{i} = M_{i-1} \cup \{ \mathbf{\check{d}}_{i,j} \mid \mathbf{\check{d}}_{i,j} \in D_{i}, \\ \nexists \mathbf{m}_{i-1,k} \in M_{i-1} : \| \mathbf{\check{d}}_{i,j} - \mathbf{m}_{i-1,k} \| < \epsilon_{\mathrm{D}} \}$$

Mobile service robot: Johnny Jackanapes

(Vice-World Champion RoboCup@Home 2008)

## **Results: Simultaneous Localization and Mapping**

2D Sparse Point Maps ( $\epsilon_D = 0.2 \text{ m}$ ), average processing time per scan: 8 ms Average processing time with Rao-Blackwellized Particle Filter and 20 particles:  $> 150 \,\mathrm{ms}$ 



The transformation  $\mathbf{T}_i$  is determined by means of the Iterative Closest Point (ICP) algorithm. Given a data set D and a model set M, the ICP iteratively searches for  $T_i$  by minimizing:

$$E(\mathbf{R}, \mathbf{t}) = \sum_{i=1}^{|M|} \sum_{j=1}^{|D|} w_{i,j} \|\mathbf{m}_i - (\mathbf{R}\mathbf{d}_j + \mathbf{t})$$
$$w_{i,j} = \begin{cases} 1, & \mathbf{m}_i \text{ corresponds to } \mathbf{d}_j \\ 0, & \text{otherwise.} \end{cases}$$
$$\mathbf{T} = \begin{pmatrix} \mathbf{R}_{ICP} & \mathbf{t}_{ICP} \\ 0 & 0 & 1 \end{pmatrix}$$

In the resulting *point maps* the robot can localize itself and other objects.

## Path Planning Procedure

- 1. Construct Voronoi diagram using Fortune's Sweepline algorithm.
- 2. Construct convex hull using Graham's

3D Sparse Point Maps (topview:  $\epsilon_D = 2 \text{ m}$ , detail views:  $\epsilon_D = 0.2 \text{ m}$ ) Average processing time per scan:  $\approx 35 \,\mathrm{ms} \,(\epsilon_{\mathrm{D}} = 0.2 \,\mathrm{m}), \approx 450 \,\mathrm{ms} \,(\epsilon_{\mathrm{D}} = 2 \,\mathrm{m})$ 



Scan algorithm.

- 3. Prune all edges of Voronoi diagram whose distance to neighboring points is smaller than the robot's width and those that lie outside of or intersect the convex hull.
- 4. Search for shortest path by means of  $A^*$ with cost function

f(n) = g(n) + h(n)

where g(n) is the length of travelled path segments and h(n) the Euclidean distance to the goal.

## **Results: Path Planning and Motion Control**

#### **Initial State:**

Robot stands in the open "FrontDoor" **Task:** Move to the "Fridge"

Computing Voronoi diagram:  $\approx 1.3 \,\mathrm{ms}$ Pruning the Voronoi diagram:  $\approx 0.8 \,\mathrm{ms}$ Planning shortest path (black):  $\approx 2.1 \,\mathrm{ms}$ 

Robot successfully followed path and reached target pose by successively applying two nonlinear motion conrollers.

